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## EXAM PAPERS PRACTICE

2.4 Graphs Mark Scheme

## Mark schemes

## Q1.

(a) All marks AO2 (analyse)

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{1}$ | 0 | 2 | 5 | 3 | 0 | 8 |
| $\mathbf{2}$ | 2 | 0 | 1 | 0 | 0 | 0 |
| $\mathbf{3}$ | 5 | 1 | 0 | 0 | 0 | 4 |
| $\mathbf{4}$ | 3 | 0 | 0 | 0 | 1 | 0 |
| $\mathbf{5}$ | 0 | 0 | 0 | 1 | 0 | 5 |
| $\mathbf{6}$ | 8 | 0 | 4 | 0 | 5 | 0 |

## Alternative answer

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{1}$ | 0 | 2 | 5 | 3 | 0 | 8 |
| $\mathbf{2}$ |  | 0 | 1 | 0 | 0 | 0 |
| $\mathbf{3}$ |  |  | 0 | 0 | 0 | 4 |
| $\mathbf{4}$ |  |  |  | 0 | 1 | 0 |
| $\mathbf{5}$ |  |  |  |  | 0 | 5 |
| $\mathbf{6}$ |  |  |  |  |  | 0 |
|  |  |  |  |  |  |  |

PRACTICE

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{1}$ | 0 |  |  |  |  |  |
| $\mathbf{2}$ | 2 | 0 |  |  |  |  |
| $\mathbf{3}$ | 5 | 1 | 0 |  |  |  |
| $\mathbf{4}$ | 3 | 0 | 0 | 0 |  |  |
| $\mathbf{5}$ | 0 | 0 | 0 | 1 | 0 |  |
| $\mathbf{6}$ | 8 | 0 | 4 | 0 | 5 | 0 |

## Mark as follows:

1 mark 0s in correct places
1 mark all other values correct
I. non-zero symbols used to denote no edge but only for showing no edge going from a node to itself
(b) All marks for AO1 (understanding)

Adjacency list appropriate when there are few edges between vertices // when graph/matrix is sparse; NE. few edges

Adjacency list appropriate when edges rarely changed;
Adjacency list appropriate when presence/absence of specific edges does not need to be tested (frequently);
A. Alternative words which describe edge, eg connection, line, arc

## Max 2

(c) Mark is for AO2 (apply)

It contains a cycle / cycles;
(d) Mark for AO1 (knowledge)

A graph where each edge has a weight/value associated with it;
(e) All marks AO2 (apply)


## Mark as follows:

I. output column

1 mark first value of A is 2
1 mark second value of A is 5 and third value is 3
1 mark fourth and subsequent values of A are $8,3,7,4,9$ with no more values after this
1 mark $D[2]$ is set to 2 and then does not change


Page 5 of 17

|  |  |  |  | D |  |  |  |  |  | P |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | Q | v | A | 1 | 2 | 3 | 4 | 5 | 6 | 1 | 2 | 3 | 4 | 5 | 6 |
| - | $\begin{gathered} 1,2 \\ 3, \\ 4,5 \\ 6 \end{gathered}$ | - | - | 20 | 20 | 20 | 20 | 20 | 20 | -1 | -1 | -1 | -1 | -1 | -1 |
|  |  |  |  | 0 |  |  |  |  |  |  |  |  |  |  |  |
| 1 | $\begin{gathered} \hline 2,3 \\ 4, \\ 5,6 \end{gathered}$ | 2 | 2 |  | 2 |  |  |  |  |  | 1 |  |  |  |  |
|  |  | 3 | 5 |  |  | 5 |  |  |  |  |  | 1 |  |  |  |
|  |  | 4 | 3 |  |  |  | 3 |  |  |  |  |  | 1 |  |  |
|  |  | 6 | 8 |  |  |  |  |  | 8 |  |  |  |  |  | 1 |
| 2 | $\begin{gathered} 3,4 \\ 5, \\ 6 \end{gathered}$ | 3 | 3 |  |  | 3 |  |  |  |  |  | 2 |  |  |  |
| 3 | $\begin{gathered} 4,5 \\ , 6 \end{gathered}$ | 6 | 7 |  |  |  |  |  | 7 |  |  |  |  |  | 3 |
| 4 | 5,6 | 5 | 4 |  |  |  |  | 4 |  |  |  |  |  | 4 |  |
| 5 | 6 | 6 | 9 |  |  |  |  |  |  |  |  |  |  |  |  |
| 6 | . |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

Max 6 marks if any errors
(f) Mark is for AO2 (analyse)

The shortest distance / time between locations/nodes 1 and 6;
NE distance / time between locations/nodes 1 and 6
R. shortest route / path
(g) All marks AO2 (analyse)

Used to store the previous node/location in the path (to this node);
Allows the path (from node/location 1 to any other node/location) to be recreated // stores the path (from node/location 1 to any other node/location);

Max 1 if not clear that the values represent the shortest path

## Alternative answer

Used to store the nodes that should be traversed;
And the order that they should be traversed;
Max 1 if not clear that the values represent the shortest path

Q2.
(a) Mark is for AO1 (understanding)

It contains a cycle / cycles;
(b) All marks AO2 (apply)

| Vertex <br> (in <br> Figure $\mathbf{1}$ | Adjacent <br> vertices |
| :---: | :--- |
| $\mathbf{1}$ | 2,3 |
| $\mathbf{2}$ | $1,3,4$ |
| $\mathbf{3}$ | $1,2,5$ |
| $\mathbf{4}$ | 2 |
| $\mathbf{5}$ | 3 |

Mark as follows:
(c) All marks AO1 (understanding)

Adjacency list appropriate when there are few edges between vertices / / when graph / matrix is sparse;
when edges rarely changed;
when presence / absence of specific edges does not need to be tested (frequently);
Max 2
A Alternative words which describe edge, eg connection, line.
(d) All marks AO2 (apply)

|  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| NoOfCats | A | B | C | 1 | 2 | 3 | 4 | 5 |



## Mark as follows:

1 mark: A is set the sequence indicated in the table;
1 mark: B is set the sequence indicated in the table;
1 mark: C is set the sequence indicated in the table;
1 mark: NoOfCats is set to 5 , Cat [1] is set to 1 ;
1 mark: Cat [2] is set to 2 and Cat [3] is set to 3;
1 mark: Cat [4] is set to 1 and Cat [5] is set to 1 ;
Info for examiner: Ignore the empty cells in the sequences - values do not need to be set in the rows indicated in the table.
(e) Mark is for AO2 (analyse)

To work out which cats will travel together to the show / /
To plan which cats will be in the van on which journey to the cat show //
To colour the vertices of a graph / /

To create a decomposition of a graph; Max 1
(f) All marks AO1 (knowledge)

1 mark (1 from): The problem can be solved / / algorithm exists for problem; But it cannot be solved in polynomial time / / but not quickly enough to be useful;
Max 2
1 mark: It takes an unreasonable amount of time; to solve;
A Too long time but $\mathbf{R}$ Long time
(g) All marks AO1 (understanding)

1 mark: Use of heuristic; algorithm that makes a guess based on experience; That provides a close-to-optimal solution / approximation; that only works in some cases; A non-optimal

Example of heuristic method eg hill-climbing / stochastic / local improvement / greedy algorithms / simulated annealing / trial and error / any reasonable example;

1 mark: Relax some of the constraints on the solution; A Solve simpler version of problem


Q3.
(a)


If a letter has been used more than once then mark it as correct in the row that it is correct in, if any.
(b) Example 1

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{1}$ | 0 | 1 | 0 | 1 | 0 | 0 |
| $\mathbf{2}$ | 1 | 0 | 1 | 0 | 1 | 0 |
| $\mathbf{3}$ | 0 | 1 | 0 | 0 | 1 | 1 |
| $\mathbf{4}$ | 1 | 0 | 0 | 0 | 1 | 0 |
| $\mathbf{5}$ | 0 | 1 | 1 | 1 | 0 | 0 |
| $\mathbf{6}$ | 0 | 0 | 1 | 0 | 0 | 0 |

## Example 2



1 mark for labelling matrix with indices running from 1 to 6 on both axis and filled only with 0 s and 1 s , or some other symbol to indicate presence / absence of edge. e.g. T / F (allow a third symbol along diagonal). Absence can be represented by an empty cell.
1 mark for correct values entered into matrix, as shown in either example above

In Example 2, the shaded portion can be in either half - some indication must be made that half of the matrix is not being used. This could just be leaving it blank, unless the candidate has also represented absence of an edge by leaving cells blank.

Allow use of a third symbol in the central diagonal to indicate it unused, as it would not make sense to use it in this example.

Accept column and row labels in any order so long as they correspond to the data i.e. do not have to be in sequence 1 to 6.

|  |  |  |  |  |  | Discovered |  |  |  |  |  | Parent |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| S | D | V | U | C | Queue | 1 | 2 | 3 | 4 | 5 | 6 | 12 | 23 | 4 | 56 | 6 Found | Output |
|  | - | 入 |  |  |  | F | F | F | F | F F | F | , | 入 | 4 | XX |  | $>$ |
| 1 | 6 |  |  |  | 1 | T |  |  | E |  |  |  |  |  |  | F |  |
| 1. | 6 | 1 | 2 |  | 2 |  | T |  | E |  |  |  | 1 |  |  | E |  |
| 1 | 6 | 1 | 4 |  | 24 |  | T |  | T |  |  |  |  | 1 |  | E |  |
| 1 | 6 | 2 | 1 |  | 4 |  | T | E | T |  |  |  |  | 1 |  | F |  |
| 1 | 6 | 2 | 3 |  | 43 |  |  | T | T |  |  |  | 12 | 1 |  | F |  |
| 1 | 6 | 2 | 5 |  | 435 |  | T | I | T | T |  |  | 2 | 12 | 2 | E |  |
| 1 | 6 | 4 | 1 |  | 35 |  | T | I | T | T |  |  | 12 | 1 |  | E |  |
| 1 | 6 | 4 | 5 |  | 35 |  | T | T | I | T |  |  | 12 | 1 |  | E |  |
| 1 | 6 | 3 | 2 |  | 5 |  | I | T | I | I |  |  | 12 | 1 | 2 | F |  |
| 1 | 6 | 3 | 5 |  | 5 |  | T | T | T | T |  |  | 12 | 1 | 2 | F |  |
| 1 | 6 | 3 | 6 |  | 56 |  | T | T | T | T | T |  | 12 | 1 | 23 | 3 T |  |
| 1 | 6 | 3 | 6 | 6 | 56 | 1 | T | T | T | T | T |  | 12 | 1 | 2 | I | 6 |
| 1 | 6 | 3 | 6 | 3 | 56 | T | I | T | T | T | T |  | 2 | 1 | 2 | I | 3 |
| 1 | 6 | 3 | 6 | 2 | 56 | T | T | T | T | T | - |  | 12 | 1 | 2 | I | 2 |
| 1 | 6 | 3 | 6 | 1 | 56 | T | T | T | T | T | $\pm$ |  | 12 | 1 | 23 | I | 1 |

1 mark for having the correct value changes in each region highlighted by a rectangle and no incorrect changes in the region. Ignore the contents of any cells that are not changed.
A Alternative indicators that clearly mean True and False.
A It is not necessary to repeat values that are already set (shown lighter in table)
A For the queue column, for cells that should only have one value in them, accept if the student has written out the value twice at both the Front and the Rear, solong as this has been done consistently throughout the table.eg " 4 " written as "4 4".
(d) So that packets / data arrive as quickly as possible / more quickly / with lower latency;
So that packets / data are delivered at lower cost / lowest cost;
So that routers do not have to process more packets / data due to
unnecessary hops being made;
NE it is quicker without clarifying what "it" refers to
MAX 1

Q4.
(a) Connected // There is a path between each pair of vertices; Undirected // No direction is associated with each edge;
Has no cycles // No (simple) circuits // No closed chains // No closed paths in which all the edges are different and all the intermediate vertices are different // No route from a vertex back to itself that doesn't use an edge more than
once or visit an intermediate vertex more than once;
A no loops
Alternative definitions:
A simple cycle is formed if any edge is added to graph;
Any two vertices can be connected by a unique simple path;
(b) No route from entrance to exit / through maze;

Maze contains a loop/circuit ;
A more than one route through maze;
Part of the maze is inaccessible / enclosed;
$\mathbf{R}$ Responses that clearly relate to a graph rather than the maze
(c)

|  | $\mathbf{1}$ | $\mathbf{2}$ | $\mathbf{3}$ | $\mathbf{4}$ | $\mathbf{5}$ | $\mathbf{6}$ | $\mathbf{7}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $\mathbf{1}$ | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| $\mathbf{2}$ | 1 | 0 | 1 | 1 | 0 | 0 | 0 |
| $\mathbf{3}$ | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| $\mathbf{4}$ | 0 | 1 | 0 | 0 | 1 | 0 | 0 |
| $\mathbf{5}$ | 0 | 0 | 0 | 1 | 0 | 1 | 1 |
| $\mathbf{6}$ | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| $\mathbf{7}$ | 0 | 0 | 0 | 0 | 1 | 0 | 0 |

(allow some symbol in the central diagonal to indicate unused)

(with the shaded portion in either half - some indication must be made that half of the matrix is not being used. This could just be leaving it blank, unless the candidate has also represented absence of an edge by leaving cells blank)

1 mark for drawing a $7 \times 7$ matrix, labelled with indices on both axis and filled only with 0s and 1s, or some other symbol to indicate presence/absence of edge. e.g. T/F. Absence can be represented by an empty cell.
1 mark for correct values entered into matrix, as shown above;
(d) (i) Routine defined in terms of itself // Routine that calls itself;

A alternative names for routine e.g. procedure, algorithm
NE repeats itself
(ii) Stores return addresses;

Stores parameters;
Stores local variables; NE temporary variables
Stores contents of registers;
A To keep track of calls to subroutines/methods etc.

## Max 1

Procedures / invocations / calls must be returned to in reverse order (of being called);
As it is a LIFO structure;
A FILO
As more than one / many return addresses / sets of values may need to be stored (at same time) // As the routine calls itself and for each call/invocation a new return address / new values must be stored;

Max 1
(e)

|  |  |  |  | Discovered |  |  |  |  |  |  |  | Completely Explored |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Call | V | U | $\begin{aligned} & \mathrm{En} \\ & \mathrm{dV} \end{aligned}$ | 1 | 2 | 3 | 4 | 5 | 6 | 6 | 7 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | F |
|  | - | - | 7 | F | F | F | F | F | F | F | F | F | F | F | F | F | F | F | F |
| DFS(1,7) | 1 | 2 | 7 | T | F | F | F | F | F |  |  |  | F | F | F | F | F | F |  |
| DFS(2,7) | 2 | 1 | 7 | T | T | F | F | F | F | F 1 |  |  | F | F | F | F | F | F |  |
|  |  | 3 | 7 | I | T | F | F | F | F |  |  |  | F | F | F | F | F | F |  |
| DFS(3,7) | 3 | 2 | 7 |  | T T | T | F | F | F |  |  |  | F | T | F | F | F | F |  |
| DFS(2,7) | 2 | 4 | 7 | I | T | T | F | F | F |  |  |  | F | T | F | F | F | F |  |
| DFS(4,7) | 4 | 2 | 7 | [ | T | T | T | F | F |  |  |  | F | T | F | F | F | F |  |
|  |  | 5 | 7 | I | T | T | 1 | F | H |  |  |  | F | T | F | F | F | F |  |
| DFS(5,7) | 5 | 4 | 7 | T | T | T | T | T |  |  |  |  | F | T | F | F | F | F |  |
|  |  | 6 | 7 | [ | T | T | 1 | T | F |  |  |  | F | T | F | F | F | F |  |
| DFS(6,7) | 6 | 5 | 7 | T | T | T | I | 1 |  |  |  |  | F | T | F | F | T | F |  |
| DFS(5,7) | 5 | 7 | 7 | [ | T | T | T | I |  |  |  |  | F | T | F | F | T | F |  |
| DFS(7,7) | 7 | 5 | 7 | I | T | T | I | I |  |  | T |  | F | T | F | F | T | T | T |
| DFS(5,7) | 5 | - | 7 | I | T 1 | T | T | 1 |  |  |  |  | F | T | F | T | T | T |  |
| DFS(4,7) | 4 | - | 7 | T | T | T | I | 1 |  |  |  |  | F | T | T | T | T | T |  |
| DFS(2,7) | 2 | - | 7 | [ | 1 | T | 1 | 1 |  |  |  |  | T | T | T | T | T | T |  |
| DFS(1,7) | 1 | - | 7 |  | T | T | T | T |  |  |  | T | T | T | T | T | T | T | T |

1 mark for having the correct values changes in each region highlighted by a rectangle and no incorrect changes in the region. Ignore the contents of any cells that are not changed.

A alternative indicators that clearly mean True and False.
A it is not necessary to repeat values that are already set (shown lighter in table)

## Q5.

(a)


1 mark for all 5 lines correctly drawn
1 mark for all 5 arrowheads pointing in correct directions
Max 1 if more than 5 lines drawn by candidate (note that dotted arrow is given in question)

A arrowheads at any position on line
(b) Adjacency matrix appropriate when there are many edges between vertices // when edges may be frequently changed // when presence/absence of specific edges needs to be tested (frequently)
Adjacency list appropriate when there are few edges between vertices // when graph is sparse // when edges rarely changed //when presence/absence of specific edges does not need to be tested (frequently)
A alternative words which describe edge e.g. connection, line
(c) Connected // There is a path between each pair of vertices; $\square$ Undirected // No direction is associated with each edge;
Has no cycles // No (simple) circuits // No closed chains // No closed paths in which all the edges are different and all the intermediate vertices are different // No route from a vertex back to itself that doesn't use an edge more than once or visit an intermediate vertex more than once;

## Alternative definitions:

Graph with no cycles, and a simple cycle is formed if any edge is added to it;; Graph which is connected, and it is not connected anymore if any edge is removed from it;;
Graph in which any two vertices can be connected by a unique simple path;;
(Note: not just adjacent vertices)
Graph which is connected and has $\mathrm{n}-1$ edges where n is no of vertices;; Graph which has no simple cycles and has $\mathrm{n}-1$ edges where n is no of vertices;
(d)


1 mark for Jack as root
1 mark for Bramble and Snowy as children of Jack
1 mark for four correct children of Bramble and Snowy
DPT if arrows drawn instead of lines
DPT if any node has more than 2 child nodes
A "mirror image" answers which are consistent.
(e) For solution with 3 arrays:

One array stores data items;
One array for left child pointers;
One array for right child pointers;
Pointers stored at same location in arrays as corresponding data item;
For solution with 1 array of records:
Record created to store data item and pointers;
One pointer to left child;
One pointer to right child;
For either of the above solutions:
Rogue value (allow example) used to indicate no child;
Variable indicates position in array(s) of root node // Root node stored at first location/start of array(s);

## If answered as diagram:

Column for data with at least three correct data items in it,
Use of rogue value for a node that does not have child;
Correct value for a start pointer variable indicating position of root node in the array (not drawn as an arrow, array indices must be labelled);
Column for left child pointers*;
Column for right child pointers*;

* $=$ To get these marks, there must be a sufficient number of pointers to demonstrate that the data structure is a representation of a binary tree, but it is not necessary for every item to be shown. Also the array indices must be shown.


## Q1.

In previous years there have been questions asking students to complete an adjacency matrix based on a diagram of a graph and most students were able to answer question (a) this year. This was the first time that an adjacency matrix for a weighted graph had been asked for and some students had clearly not seen this type of question before and only included an indicator that there was an edge between two nodes rather than the weight of the edge between the two nodes; this meant they only got one of the two marks available for this question.

Questions (b)-(d) were about graph theory. Question (c) was well-answered with students identifying that it was not a tree because there were cycles. The most common incorrect answer was to say that it wasn't a tree because the edges have weights associated with them. Question (d) was also well-answered. Answers to (b) often showed that students were not as familiar with adjacency lists as they are with adjacency matrices.

For question (e) students had to complete a trace of Djikstra's Algorithm. This topic was not on the previous A-level specification and was often poorly answered suggesting many students had not tried to complete a trace for the algorithm before. For question (f) many students gave an answer that explained the point of Djikstra's Algorithm (find the shortest route from a node to each other node) rather than what the specific output in the algorithm given in the question would be (the distance of the shortest route from node 1 to node 6 ).

## Q3.

This question was about graphs and graph traversal in the context of a computer network.
Part (a) required students to identify graphs with specific properties. Almost all students were able to achieve one of the two marks and just over a third achieved both. Part (b) was also well tackled, with over $90 \%$ of students achieving both marks for correctly completing the adjacency matrix.

For part (c) students had to trace an algorithm for performing a breadth first search of a graph. This was well tackled, with over half of students achieving at least four of the six marks. The most commonly made mistakes were to incorrectly record the queue contents and to update the Parent array incorrectly.

For part (d) students needed to explain why finding the shortest route was a useful thing to do in the context of the question. Just over half of students correctly explained that this would allow data to be transmitted more quickly or would reduce congestion. Some students missed out on achieving a mark by giving responses that were not in context.

Q4.
Part (a): Two thirds of students were able to identify one property that a graph must have to be a tree. A small number confused a tree with a rooted tree and made assertions such as that a tree must have a root, which is incorrect.

Part (b): This question part tested students' understanding of the method being used to represent a maze as a graph. The majority of students correctly identified a feature of the maze that would stop its graph being a tree. The most commonly seen correct response identified that there could be a loop in the maze. Other possibilities included that part of the maze could be inaccessible or that part of the maze might only be traversable in one direction. Some students failed to achieve the mark because they re-answered part (a),
discussing a feature of a graph that would stop it being a tree, rather than a feature of a maze.

Part (c): Students were asked to represent the graph of the maze as an adjacency matrix. Three quarters of students scored both marks for this question part. Responses where symbols other than 0s and 1s were used in the matrix were accepted, as long as they could be viewed as an accurate representation of the graph.

Part (d)(i): The vast majority of students were able to identify that a recursive routine would call itself. A small number asserted that a recursive routine would repeat itself, which was not considered to be enough for a mark as this could equally have been a description of iteration.

Part (d)(ii): Most students scored some marks for this question part, but less than a fifth achieved both. The most widely understood point was that the data would need to be removed from the stack in the reverse of the order that it was put onto it so that the recursion could be unwound. Less well understood was the types of data that would be stored, such as return addresses and local variables.

Part (e): Most students achieved some marks on this question part and around a quarter achieved all five for a fully complete trace. The most commonly made mistake was to update, incorrectly, the Completely Explored array as the recursive calls were made, as opposed to when the recursion unwound.

Q5.
Part (a): The use of the adjacency matrix was clearly well understood with all but a few candidates achieving full marks.

Part (b): There were some good responses to this question part, but also quite a lot of confused answers. An adjacency matrix is more appropriate when there are many edges in a graph, or if these edges need to be checked or updated frequently. An adjacency list is appropriate for graphs with few edges (sparse graphs) or where the edges are not checked / updated frequently.
Neither the number of vertices in a graph nor the available memory would influence the choice.
There was confusion over the use of terminology with some candidates apparently using the term vertex to mean edge.

Part (c): This question part was poorly answered, with only a third of candidates scoring any marks. A tree is a graph that is connected, undirected and has no cycles. Some candidates gave responses that referred only to specific types of tree, either rooted trees or binary trees.
These responses did not gain credit.
Part (d): The vast majority of candidates knew how to construct a binary search tree. The most common cause of error appeared to be candidates forgetting the order of the letters in the alphabet rather than forgetting the principles that should be used to construct the tree. A small number of candidates mistakenly drew arrows instead of lines between nodes.

Part (e): There were some good responses to this question part but many were disappointing and a surprising number of candidates did not write a response at all. Many candidates who did answer chose to use a diagram to illustrate their response which was quite acceptable, so long as the diagram included enough detail to make clear that it was a representation of a binary tree.

